

SkiNet

A Petri Net Generation Tool for the Verification of Skillset-based Autonomous Systems

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Workshop on Formal Methods for Autonomous Systems

Presentation plan

- 1 Context
- 2 Skillset
- 3 Translation
- 4 Verification
- 5 Conclusion
- 6 Background

Skill-based architecture?

This work verifies skill-based specifications of robots, as introduced by Albore et al. [1].

Skill-based software architecture:

Pros

- Simpler specification
- Easier verification
- Ideal for automation
- Doesn't need a high experience in robotics
- Controller code is generated from the model and follows the same semantic

Cons

- Requires a solid middleware
- Hardware issues are either overlooked or unpredictable
- Scalability issues for verification with increasing complexity

[1] Alexandre Albore, David Doose, Christophe Grand, Charles Lesire and Augustin Manecy (2021): Skill-Based Architecture Development for Online Mission Reconfiguration and Failure Management. In: 2021 IEEE/ACM 3rd International Workshop on Robotics Software Engineering (RoSE), IEEE, pp. 47–54, doi: 10.1109/RoSE52553.2021.00015.

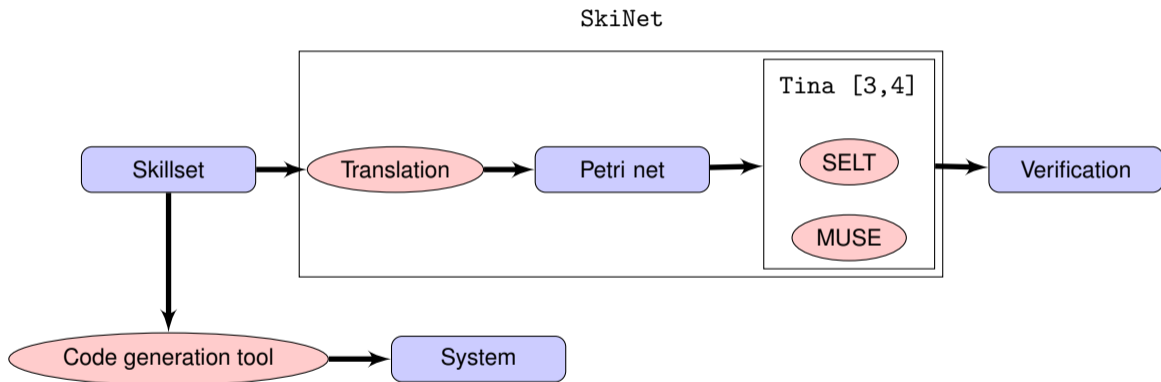
How much do you trust your robot?

Objective of SkiNet: improve the **trustability** of autonomous systems designed with **skill-based architectures** (Albore et al. [1]).
Using **a-priori Verification**:

- How can we check whether our **high-level model** of the system is not flawed?
- Are there **undesired or blocking states**?
- Can it be polyvalent and perform a wide range of **missions**?
- Can we make the verification tool **easy to use and specify**?



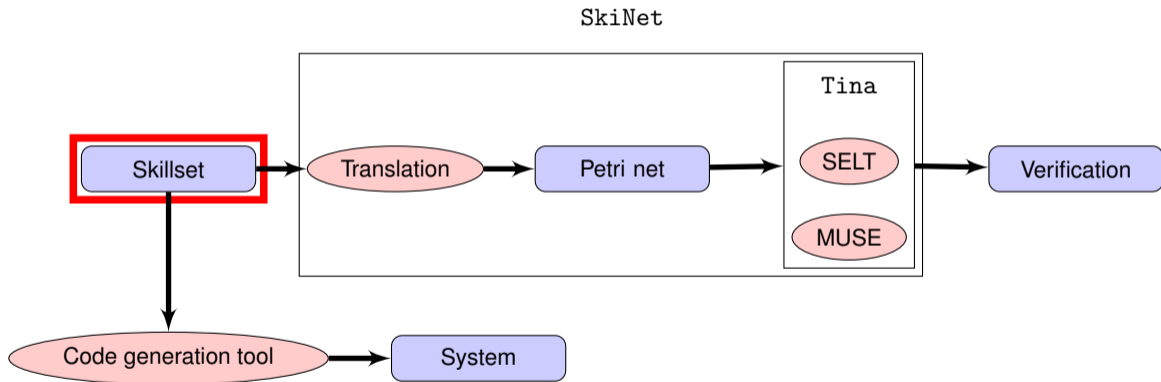
SkiNet - Tool overview



[3] B. Berthomieu, P.-O. Ribet and F. Vernadat (2004): The tool TINA – Construction of abstract state spaces for petri nets and time petri nets. International Journal of Production Research 42, pp. 2741–2756, doi:10.1080/00207540412331312688.

[4] Bernard Berthomieu, François Vernadat and Silvano dal Zilio (2004): The TINA toolbox Home Page - Time Petri Net Analyzer - by LAAS/CNRS. Available at <https://projects.laas.fr/tina/home.php>

SkiNet - Skillset



Example system - Spot[®] skillset architecture

Skillset resources

- power_status: PowerOff PowerOn
- lease_status: AutoMode ManualMode
- control_mode: Idle Busy

Skills

- init_power
- safe_poweroff
- go_to



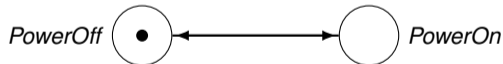
Resources

```
skillset spot_fmas {  
  
  resource {  
    power_status {  
      initial PowerOff  
      PowerOff -> PowerOn  
      PowerOn -> PowerOff  
    }  
    lease_status {  
      initial AutoMode  
      AutoMode -> ManualMode  
      ManualMode -> AutoMode  
    }  
    control_mode {  
      initial Idle  
      Idle -> Busy  
      Busy -> Idle  
    }  
  }  
}
```

Resources \mathcal{R} = finite-state machines (FMSs)

- **Internal elements:** battery level, stance, power status, control mode, mutex, sensor data...
- **External elements:** operator control, environment objects, mission elements...

Ex. power_status:

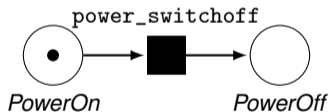


Events - Resource guards and effects

```
event {  
  toauto_frommanual {  
    guard lease_status == ManualMode  
    lease_status -> AutoMode  
  }  
  tomanual_fromauto {  
    guard lease_status == AutoMode  
    lease_status -> ManualMode  
  }  
  power_switchoff {  
    guard power_status == PowerOn  
    power_status -> PowerOff  
  }  
  power_switchon {  
    guard power_status == PowerOff  
    power_status -> PowerOn  
  }  
}
```

Events \mathcal{V} = **uncontrollable elements** that can affect the state of resources.

- **Resource guard** = logical formula
 $\phi : \{S^r, r \in \mathcal{R}\} \rightarrow \{True, False\}$
- **Effect** $\epsilon = (r, S_i^r)$ of a resource and its **next state**.

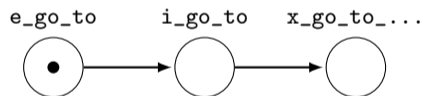


Skills

Skills \mathcal{S} = elementary actions FSMs:

- An id.
- **States:** idle e_s , running i_s and ended $x_{s,k}$
- ...

```
skill go_to {  
    ...
```

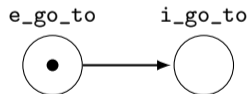


Skills

Skills \mathcal{S} = elementary actions FSMs:

- ...
- **Preconditions** (start guards). Failure effects can happen when skill starts but guards aren't satisfied.
- **Start** effects if all preconditions are satisfied.
- ...

```
skill go_to {  
  precondition {  
    is_powered { guard (power_status==PowerOn) }  
    is_auto { guard (lease_status==AutoMode) }  
    is_idle { guard (control_mode==Idle) }  
  }  
  start control_mode -> Busy  
  ...  
}
```



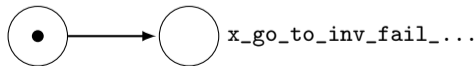
Skills

Skills \mathcal{S} = elementary actions:

- ...
- **Invariants** (execution guards). Skill ends immediately upon failure, with failure effects.
- ...

```
skill go_to {  
  ...  
  invariant {  
    is_powered {  
      guard (power_status==PowerOn)  
      failure control_mode -> Idle  
    }  
    is_auto {  
      guard (lease_status==AutoMode)  
      failure control_mode -> Idle  
    }  
  }  
  ...  
}
```

i_go_to

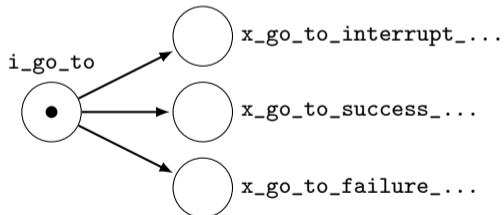


Skills

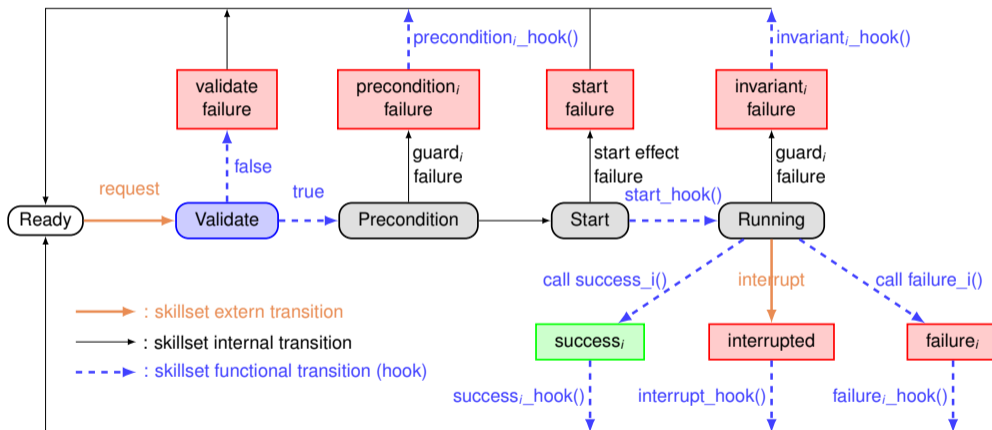
Skills S = elementary actions:

- ...
- **Interrupt** (external operation).
- **Success** and **Failure** (internal operations).

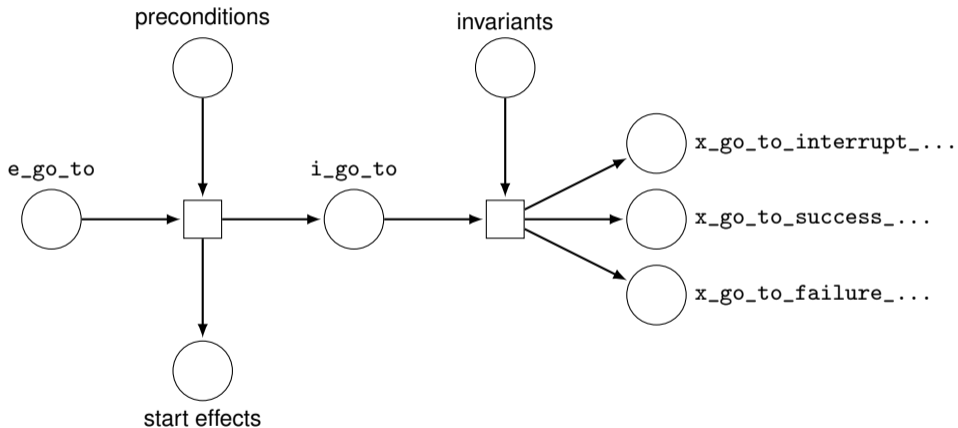
```
skill go_to {  
    ...  
    interrupt control_mode -> Idle  
    success is_arrived control_mode -> Idle  
    failure failed_reach control_mode -> Idle  
}
```



Skills

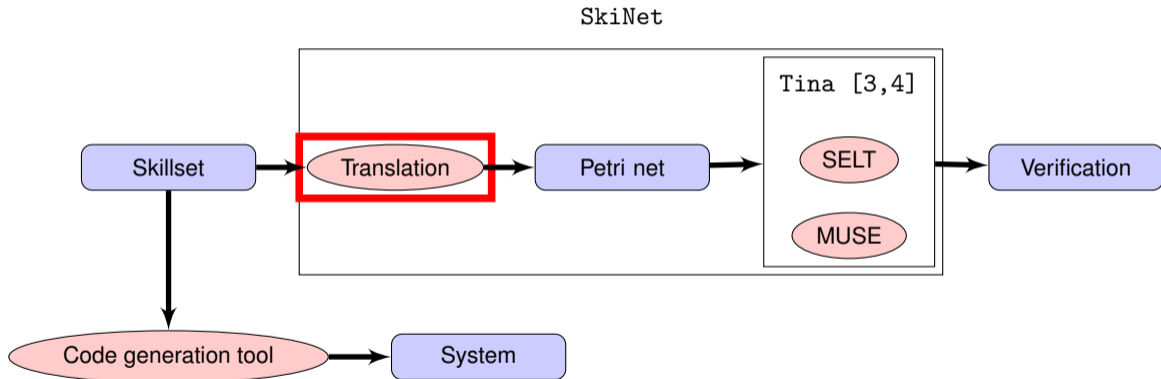


Goal for the generated Petri net



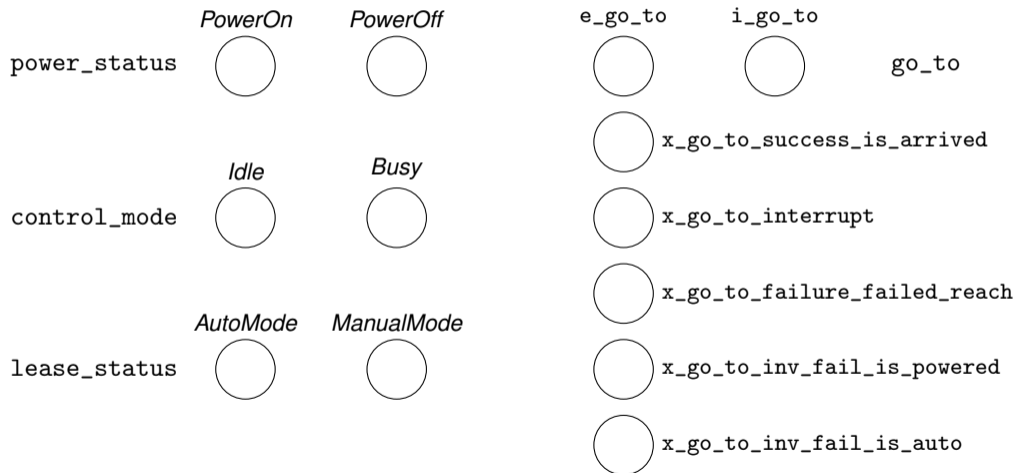
SkiNet - Translation

Petri nets = ideal for modelling **concurrent state-machines** with **shared resources**.



Translation - Places generation

Places $P = P_r$ resource state places $\cup P_s$ skills execution state



Translation - Transitions generation

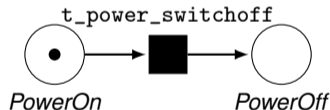
We introduce **skillset transitions** $\tau = (\phi, \mathcal{E}, \sigma)$ to generate the transitions $T = \bigcup T_\tau$.

Skillset transition τ :

- guard ϕ
- effects \mathcal{E}
- state change σ

$$\begin{aligned} \tau_{\text{power_switchoff}} = \{ & \phi = (\text{power_status} == \text{PowerOn}), \\ & \mathcal{E} = ((\text{power_status}, \text{PowerOff})), \\ & \sigma = \emptyset \} \end{aligned} \quad (1)$$

```
event power_switchoff {  
  guard power_status == PowerOn  
  power_status -> PowerOff  
}
```

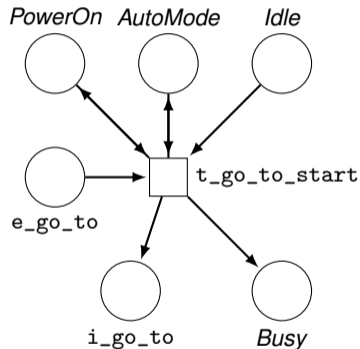


Translation - Transitions generation

```
skill go_to {
  precondition {
    is_powered { guard (power_status==PowerOn) }
    is_auto { guard (lease_status==AutoMode) }
    is_idle { guard (control_mode==Idle) }
  }
  start control_mode -> Busy
  ...
}
```

If a resource r is **evaluated** by ϕ but **not affected** by \mathcal{E} ?
→ **Reading arc**

$$\begin{aligned}
 \tau_{go_to_start} = \{ & \phi = (\text{power_status} == \text{PowerOn} \\
 & \wedge \text{lease_status} == \text{AutoMode} \\
 & \wedge \text{control_mode} == \text{Idle}), \\
 \mathcal{E} = & ((\text{control_mode}, \text{Busy})), \\
 \sigma = & e_{go_to} \rightarrow i_{go_to} \}
 \end{aligned}
 \tag{2}$$

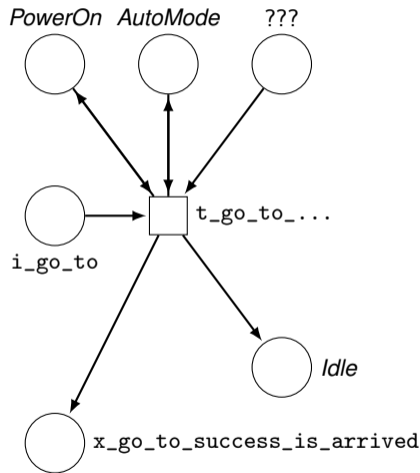


Translation - Transitions generation

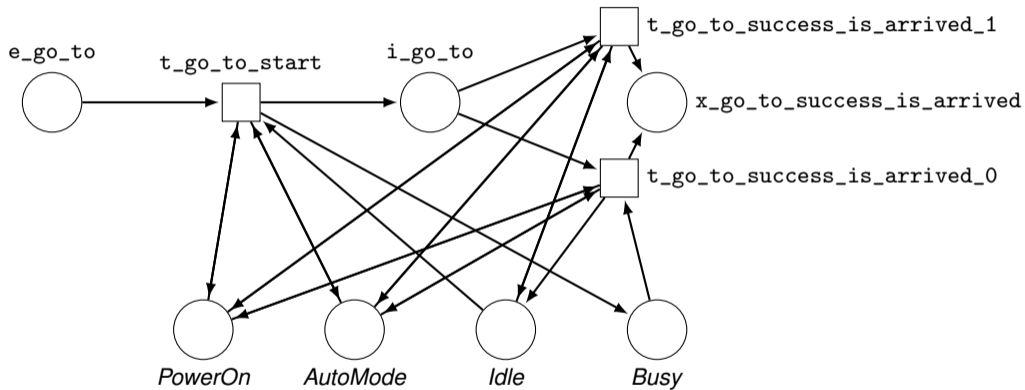
```

skill go_to {
  ...
  invariant {
    is_powered {
      guard (power_status==PowerOn)
      failure control_mode -> Idle
    }
    is_auto {
      guard (lease_status==AutoMode)
      failure control_mode -> Idle
    }
  }
  interrupt control_mode -> Idle
  success is_arrived control_mode -> Idle
  failure failed_reach control_mode -> Idle
}
  
```

If a resource r is **affected** by \mathcal{E} but **not guarded** by ϕ ?
 → **Need to anticipate the possible origin states!**

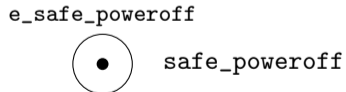
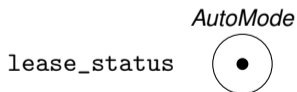
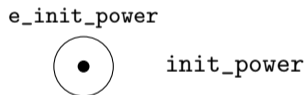
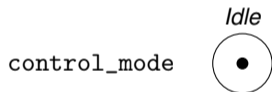
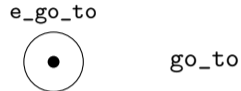
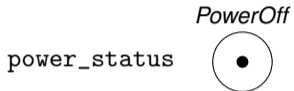


Translation - go_to start + success



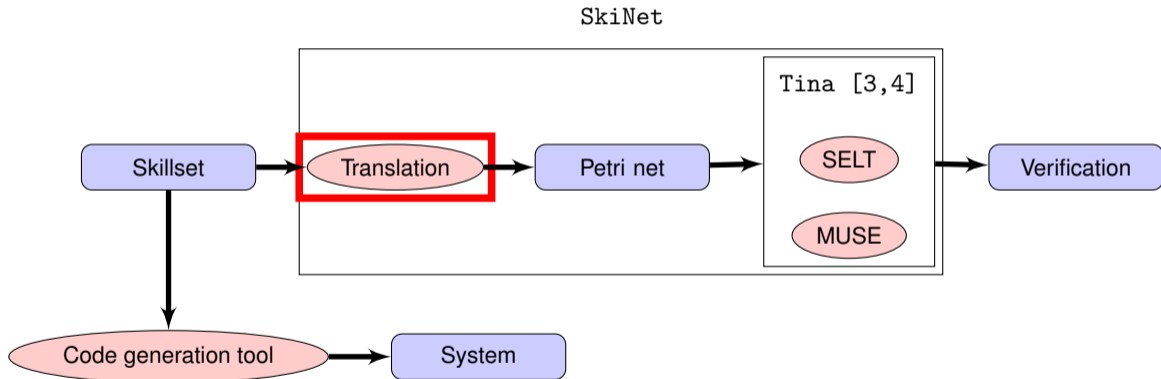
Initial marking

Resources in their **initial state**, skills are **idle**:



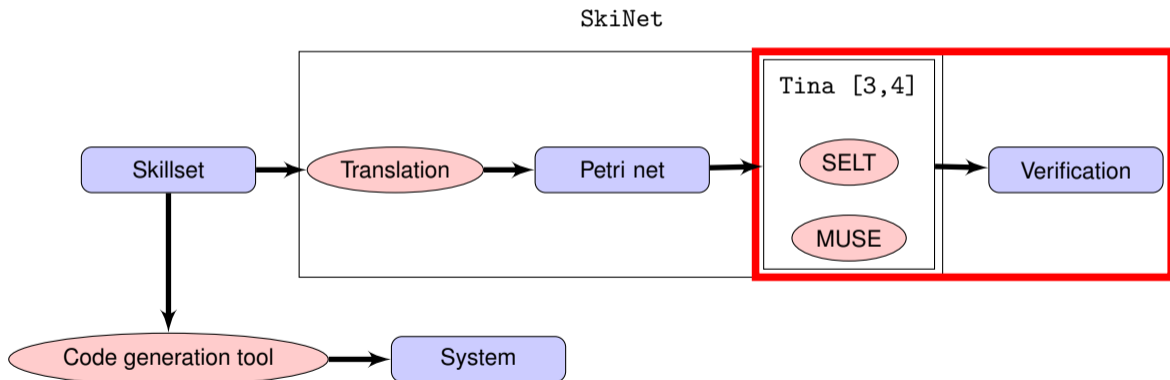
SkiNet - Translation

Generation is **formally verified** in the paper and compliant with skillset semantics.



SkiNet - Verification

After generating the **reachability graph** of the generated net, we can perform **model-checking**.



SELT - LTL Model-checker

Deadlocks → can the skillset block completely?

$$A \neg \text{dead} \tag{3}$$

Dead transitions → are there unnecessary elements?

$$\forall t \in T : A \neg t \tag{4}$$

SELT - LTL Model-checker

Check the FSM property of resources and skills (only state active at a time):

Resource invariants:

$$\forall r \in \mathcal{R} : A \left(\sum_{s_i^r \in \mathcal{S}^r} p_i^r = 1 \right) \quad (5)$$

Skill invariants:

$$\forall s \in \mathcal{S} : A \left(p_e^s + p_i^s + \sum_k p_{x,k}^s = 1 \right) \quad (6)$$

In general, check if no states are stacking:

1-safeness

$$\forall p \in P : A \neg (p \geq 2) \quad (7)$$

MUSE - mu-calculus CTL model-checker

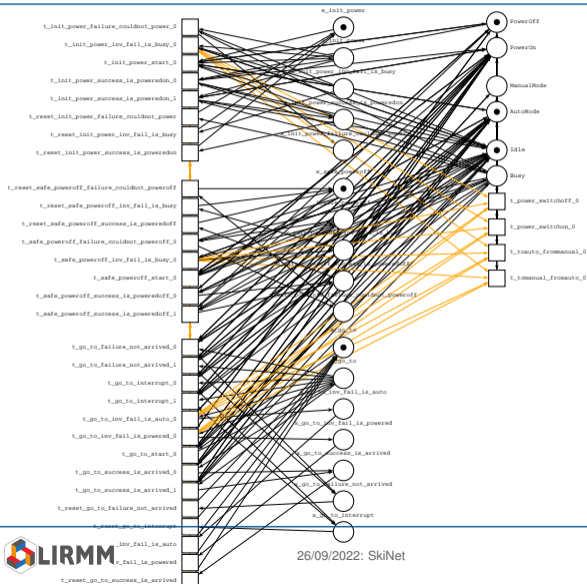
Skill liveness → Can a skill suddenly become inoperable forever?

$$\neg AFEF p_i^s, s \in \mathcal{S} \quad (8)$$

For all skills:

$$\neg AFEF \left(\sum_{s \in \mathcal{S}} p_i^s \right) \quad (9)$$

Generated net - Spot example



Verification - Spot example

```
Analysis of the skillset net.
Loading .ktz file, this can take a few seconds...
Selt version 3.7.0 -- 01/19/22 -- LAAS/CNRS
ktz loaded, 375 states, 1648 transitions
Muse version 3.7.0 -- 01/11/22 -- LAAS/CNRS
ktz loaded, 375 states, 1648 transitions
Loading successful!
Basic check: [all/dead/live/safe/skill/deadset/quit]
No deadlock: [] -dead
TRUE
Transition is never fired: [] - t
7 dead transitions found:
t_init_power_inv_fail_is_busy_0, t_reset_init_power_inv_fail_is_busy, t_reset_safe_poweroff_inv_fail_is_busy, t_safe_poweroff_inv_fail_is_busy_0, t_go_to_failure_not_arrived_1, t_go_to_interrupt_1, t_go_to_success_is_arrived_1
No more than one token per resource and skills: [] (p1 + p2 + ...) = 1
power_status
TRUE
lease_status
TRUE
control_mode
TRUE
init_power
TRUE
safe_poweroff
TRUE
go_to
TRUE
No more than one token per place: [] -( p1 >= 2 \\/ p2 >= 2 \\/ ... )
TRUE
```

Verification - Spot example

```
Skill will always eventually die: -AF EF i_skill
--init_power--
FALSE, PATH TO DEAD STATE IS:
0 -> t_power_switchon_0
2 -> t_go_to_start_0
8 -> t_power_switchoff_0
Final state: 21

--safe_poweroff--
FALSE, PATH TO DEAD STATE IS:
Identical as previous (21).

--go_to--
FALSE, PATH TO DEAD STATE IS:
0 -> t_power_switchon_0
2 -> t_go_to_start_0
8 -> t_power_switchoff_0
21 -> t_go_to_inv_fail_is_powered_0
Final state: 49

Skillset will always eventually die: -AF EF (i_sk1 /\ i_sk2 /\ ...)
FALSE, PATH TO DEAD STATE IS:
0 -> t_power_switchon_0
2 -> t_go_to_start_0
8 -> t_power_switchoff_0
21 -> t_go_to_inv_fail_is_powered_0

Basic check: [all/dead/live/tokens/skill/deadset/quit]
```

Conclusion & Future work

- We developed a tool that can **verify the global behavior** of a system based on a **skillset model**.
- We can now **test and improve the model** before implementation until we're satisfied.
- However, **scalability issues** can arise if the model is too large...

If scalability issues arise, we are developing a **runtime verification** extension of the tool that generate **rolling partial state-spaces** of the Petri net model and checks the **satisfaction of specifications**.

We also work on a **planification** extension to allow users to verify and send **high-level goals** to the system, using **sensor data**.

Git repository for the tool: <https://gitlab.com/onera-robot-skills/skinet-release>

Petri net

Why Petri nets?

Petri nets are very intuitive to use when it comes to **modelling concurrent finite state-machines and shared resources**. They are a simple mathematical tool and a lot of tools to manipulate them are available online.

A **Petri net** $\langle N, m_0 \rangle$ is a tuple $N = (P, T, F)$ and an initial marking m_0 , where:

- P and T are two non-empty, disjoint and finite sets of **places and transitions**, respectively.
- $F \subseteq (P \times T) \cup (T \times P)$ a set of **directed arcs**.
- $m_0 \in M$ the **initial distribution of tokens**, called the initial marking of the net, and $M = \{m_0, m_1, \dots, m_n\}$ the set of **all possible markings** of N .

This definition is **extended** with $N = (P, T, F, \succ)$, where \succ is the **priority relation**, represented by directed arcs between transitions, with the source transition having a higher priority.

For a transition $t \in T$, we note ${}^\bullet t$ and t^\bullet its **input and output places** respectively.

The **marking of a place** p is noted $m[p]$. The **firing of an enabled transition** $t \in T$, with $\forall p \in {}^\bullet t, m[p] \geq 1$, leads to a new marking, or reachable state, m' . All the **input places of t loose a token**, i.e. $\forall p \in {}^\bullet t, m'[p] = m[p] - 1$, and all the **output places gain one token**, i.e. $\forall p \in t^\bullet, m'[p] = m[p] + 1$.